

Announcements

- 1) Math Club talk on SVD,
Thursday, 4:00, CB 2062
- 2) Change in coding question
on HW #5 - "One" for
loop changed to "two"
(problem 4)
- 3) Scholarships! Apply online
through university's system.

If you are a masters student in applied math, apply even if you aren't full-time. These are Math department scholarships, deadline Monday April 7th.

Perturbing A

2×1 case: why is this
so much more annoying
than perturbing b ?

Given $A = \begin{bmatrix} a_1 \\ a_2 \end{bmatrix}$,

perturb to

$$\tilde{A} = \begin{bmatrix} a_1 + \delta a_1 \\ a_2 + \delta a_2 \end{bmatrix}.$$

When solving least squares using the normal equations,

we have to calculate

the pseudoinverse of A

or \tilde{A} .

We start with

$$A^*A = [\bar{a}_1 \quad \bar{a}_2] \begin{bmatrix} a_1 \\ a_2 \end{bmatrix}$$

$$= |a_1|^2 + |a_2|^2$$

$$= \left\| \begin{bmatrix} a_1 \\ a_2 \end{bmatrix} \right\|_2^2$$

(as a vector in \mathbb{C}^2)

We then compute

$$A^+ = (A^* A)^{-1} A^*$$

$$= \frac{1}{|a_1|^2 + |a_2|^2} A^*$$

$$= \begin{bmatrix} \frac{\bar{a}_1}{|a_1|^2 + |a_2|^2} & \frac{\bar{a}_2}{|a_1|^2 + |a_2|^2} \end{bmatrix}$$

Writing $\tilde{A} = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$, we get

$$(\tilde{A})^+ = \begin{bmatrix} \overline{b_1} & \overline{b_2} \\ |b_1|^2 + |b_2|^2 & |b_1|^2 + |b_2|^2 \end{bmatrix}$$

where $b_1 = a_1 + \delta a_1$,

$$b_2 = a_2 + \delta a_2.$$

We see

$$|b_1|^2 + |b_2|^2$$

$$= (a_1 + \delta a_1) (\bar{a}_1 + \overline{\delta a_1})$$

$$+ (a_2 + \delta a_2) (\bar{a}_2 + \overline{\delta a_2})$$

$$= |a_1|^2 + |a_2|^2 + 2\operatorname{Re}(a_1 \overline{\delta a_1})$$

$$+ |\delta a_1|^2 + |\delta a_2|^2$$

quadratic terms

To calculate $(\tilde{A})^\dagger$,

we have to divide

$(\tilde{A})^*$ by this quantity,

which introduces quadratic

errors instead of linear

ones.

Q': How do we keep track of these errors?

Perturbing A

Think of what perturbing

A does to the range :

recall that the image
of the unit sphere in \mathbb{C}^n
is a "hyperellipse"
in $\text{ran}(A)$.

Solving for $y = Ax = AA^+ b$.

Picture

$\text{Ran}(A)$

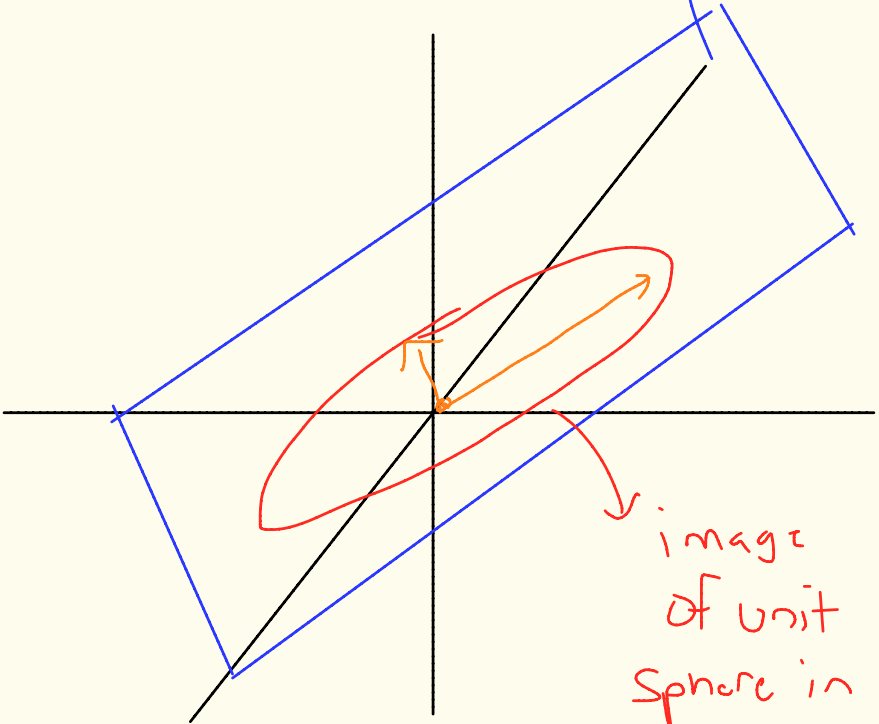


image
of unit
sphere in
 \mathbb{R}^n .

Perturb the range by a vector orthogonal to $\text{Ran}(A)$. The directions are determined by the "semi-axes" of the image of the unit sphere in \mathbb{C}^n , which are exactly the right singular vectors of A .

The smallest singular value gives you a vector with "maximum tilt", then all you care about is the angle you tilt by.

You bound the angle to bound the condition number.

Stability for Least Squares

Pick your algorithm!

But first recall

Theorem: Let $f: X \rightarrow Y$ be a problem and let \tilde{f} be a backwards stable algorithm

$\tilde{f}: X \rightarrow Y$ for f and some

$\epsilon_{\text{machine}}$. Then for all $x \in X$,

$$\frac{\|f(x) - \tilde{f}(x)\|}{\|f(x)\|} = O(K(x) \epsilon_{\text{machine}})$$

So you should expect
to lose as many
digits as the condition
number.

Using the example
in lecture 19 in
the text, we get
the following condition
numbers:

	y	x
b	1	1.1×10^5
A	2.3×10^{10}	3.2×10^{10}

Best: QR Factorization
via Householder

Algorithm: (11.2)

- 1) Compute the reduced QR decomposition of A
- 2) Compute $\hat{Q}^* b$
- 3) Solve $\hat{R} x = \hat{Q}^* b$
for x via back-substitution.

Next best: SVD

Algorithm: (11.3)

1) Compute reduced svd
of A .

2) Compute \hat{U}^* b

3) Solve $\sum w = \hat{U}^* b$
via back-substitution

4) Set $x = v w$.

algorithm?

Worst : QR Factorization via
Gram-Schmidt

Algorithm : (11.2)

Same as before, except
use Gram-Schmidt
instead of Householder
to get the reduced QR
decomposition.